AUTOMATED SYSTEM ARCHITECTURE FOR CONTAINER- BASED AND HYPERVISOR-BASED VIRTUALIZATION

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ABSTRACT

Virtualization allows multiple operating systems and applications to be executed on the same physical server concurrently. Recently, two popular virtualization platforms, namely container-based and hypervisor-based were adapted into most data centers to support cloud services. With the increase in various types of scientific workflow applications in the cloud, low-overhead virtualization techniques are becoming indispensable. However, to deploy the workflow tasks to a suitable virtualization platform in the cloud is a challenge. It requires intimate knowledge of ever-changing workflow tasks at any given moment. This research proposed an automated system architecture that can choose the best virtualization platform to execute workflow tasks. A benchmark performance evaluation was conducted on various workflow tasks running on container-based and hypervisor-based virtualization. Several tools were used to measure the metric, such as central processing unit (CPU), memory, input and output (I/O). Based on the benchmark, a system architecture was created to automate the virtualization platform selection. The results showed that the proposed architecture minimized the workflows' total execution time.

ABSTRAK

Virtualisasi membolehkan sistem operasi dan aplikasi berganda yang akan dilaksanakan pada pelayan fizikal yang sama secara serentak. Baru-baru ini, dua platform virtualisasi popular, iaitu berasaskan kontena dan berasaskan hypervisor telah disesuaikan kepada kebanyakan pusat data untuk menyokong perkhidmatan awan. Dengan peningkatan pelbagai jenis aplikasi alur kerja saintifik di awan, teknik maya kos rendah adalah sangat diperlukan. Walau bagaimanapun, ia merupakan satu cabaran untuk menggunakan tugas aliran kerja kepada platform virtualisasi yang sesuai di awan. Oleh kerana tugas aliran kerja sering berubah, ia memerlukan pengetahuan mendalam. Penyelidikan ini mencadangkan suatu senibina sistem automatik yang dapat memilih platform virtualisasi terbaik untuk melaksanakan tugas aliran kerja. Penilaian prestasi penanda aras telah dijalankan terhadap pelbagai tugas aliran kerja yang dijalankan pada virtualisasi berasaskan kontena dan berasaskan hypervisor. Beberapa peralatan digunakan untuk mengukur metrik seperti unit pemprosesan, memori, input dan output (I/O). Berdasarkan penanda aras, senibina sistem telah diwujudkan untuk mengautomasikan pemilihan platform virtualisasi. Keputusan menunjukkan bahawa senibina yang dicadangkan meminimumkan jumlah masa pelaksanaan aliran kerja.

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LIST OF SYMBOLS AND ABBREVIATIONS

- API : Application Programming Interface
- CLI : Command Line Interface
- CPU : Central Processing Unit
- DC : Data Centre
- GUI : Graphical User Interface
- I/O : Input and Output
- LXC : Linux Container
- LXD : Linux Docker
- OS : Operating System
- QoS : Quality of Services
- REST : Representational State Transfer
- SLA : Service Level Agreement
- VMs : Virtual Machines

CHAPTER 1: INTRODUCTION

This chapter begins with a background study on the container-based and hypervisorbased virtualization. Then, the challenges of the virtualizations are presented, leading to the problem statements. Subsequently, the research objectives and scopes are subsequently stated. Finally, the dissertation organisation is presented at the end of the chapter.

1.1 Background

Virtualization is a software that isolates physical infrastructures to create numerous dedicated resources. Virtualization makes it doable to run several operating systems and applications on similar server. The benefit of virtualizations is on the server-side, where the virtualization itself reduces maintenance and energy costs as well as the number of physical servers. By sharing resources in the physical servers, virtualization becomes the fundamental technology that powers Cloud computing. Figure 1.1 shows the adoption of virtualization in most data centres throughout the regions.

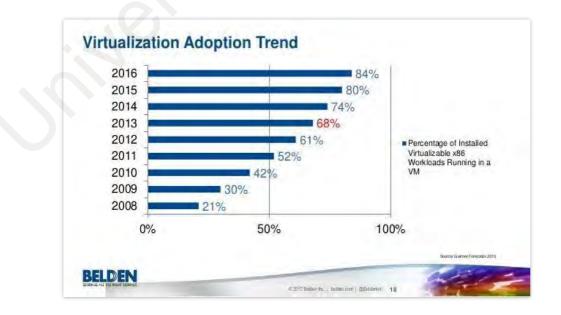


Figure 1.1. Virtualization adoption trend (Gartner, 2012)

Virtualization plays a vital role in supporting Cloud services from resource provisioning to isolating the resources. There are two popular types of virtualization platform, namely the hypervisor-based and the container-based. The hypervisor-based virtualization can be divided into the bare-metal hypervisor that is usually installed straight forwardly onto the server, and the hosted hypervisor that requires a host operating system (OS).

Figure 1.2 depicts the hypervisor-based virtualization. The bottom layer is where the physical server or the hardware is located. This layer consists of the CPU, memory and network card that are attached into one physical server. The second layer is the virtual machine monitor, namely the hypervisor. OS is isolated by the hypervisor from the hardware by taking the responsibility of permitting every running OS time with the underlying hardware. Each OS is controlled by the hypervisor called the guest OS, and the hypervisors operating system is called the host OS. This layer can handle different types of guest OS.

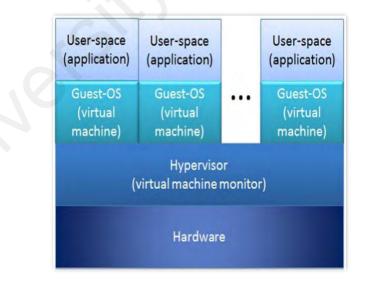


Figure 1.2. Hypervisor-based virtualization

The third layer is the virtual machine OS that is managed by the hypervisor itself. All these virtual machines have their own services and libraries. The applications are installed on each of the respective virtual machine. This virtualization allows users to virtualize many operating systems in one piece of hardware.

To reduce the performance overhead of hypervisor-based virtualization, professionals and specialists as of late began advancing an option and lightweight plans, namely the container-based virtualization. Figure 1.3 shows the container-based virtualization.

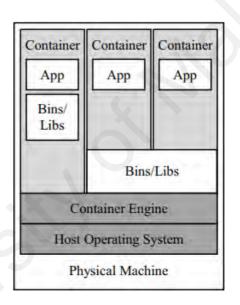


Figure 1.3. Container-based Virtualization

Hypervisor has a different kind of hardware-level solutions than the container. The container realizes within the OS level and protect their contained application by utilizing them in the slices on the host (Bernstein, 2014). The containerization is an OS-level virtualizations method of deploying and running distributed applications without launching the entire VM for each application. Instead of multiple isolated systems, the

container is executed on a single control host and accessing a single kernel. There are a few types of container-based virtualization such as Docker, LXD and lmctfy.

Container that is introduced by Docker is a free-source platform for user to develop their software. The benefit of the docker is it can group the application to become a package in "container" allowing them to be moveable among any system running the Linux OS. Docker splits the applications from the infrastructure to speed up the time taken for software deployment. Docker architecture consist of Docker engine which responsible for client-server application. There are some major components in the engine, namely daemon process which is a type of zero-downtime running program, the Representational State Transfer (REST) application programming interface (API) which specifies screens that enable programs to talk to the daemon and send command on what to do. A command line interface is important to configure, build and maintain the Docker environment. The command line interface (CLI) enables Docker REST API to manage or communicate with the Docker daemon through scripting or one-to-one CLI commands. Besides, there are some other Docker applications that use the underlying API and CLI which are shown in Figure 1.4. The daemon manages Docker objects, such as image, data volumes, network and container.

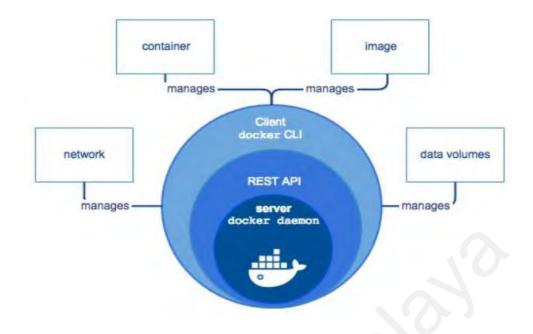


Figure 1.4. Major component in Docker engine

Hypervisor-based	Container-based
VM contain a full operating system with	Containers are executed with the
their own management of memory installed	compute engine. Containers comes
with the mixed overhead of device drivers	with more portable and less in size
virtually.	than VM and enables fast start-up
	with optimum execution.
Every virtual machine has their own	Single kernel can be shared in a
services and have their own personal	container and at the same time share
libraries	the same application libraries.
Fully loaded with resources and heavy and	Faster and less heavy resources and
consume time to create and launch VM	very reliable on fast start-up and
	launching

 Table 1.1 Differences between hypervisor-based and container-based virtualizations

Table 1.1 shows the gap between running applications in hypervisor-based and container-based virtualizations. Based on the comparison, the container is lighter and faster as compared to hypervisor-based virtualization. From the performance perspective, containers approach is better and reliable in term of start-up time and the time consumed to deploy applications on environment.

By utilizing virtualization technology, sharing resources among numerous users decrease the resource consumed. By renting the infrastructure from public Cloud, it will be a cost effective for user and the decrease time to conserve the highly priced infrastructure. Therefore, cluster and grid computing technology will have an option which is used for massive production application or scientific workflow.

Scientific workflows consist of multiple tasks that are used to conduct an experiment. Scientific workflows can vary from simple to complex tasks, and from just a few tasks to millions of tasks. Parallel and serial are the type of task that can be use in scientific workflows. These tasks have distinctive characteristics, namely compute intensive, memory intensive, I/O intensive and data intensive.

Compute intensive task is the task that consume CPU of its host and performs computationally intensive work that does not fit comfortably into the traditional platform. This compute intensive task needs an asynchronous submission and most likely can run in extended periods of time. There are a few types of compute intensive workflow tasks. These includes executing a search for a prime number that involve many big integer divisions, which calculate a large factorial like 2000! and involve many big integer multiplications.

I/O intensive task is basically the task that reads or writes a large amount of data. The performance of such task is depending on the speed of the devices or the platform that being used. The example of I/O intensive task is when the task is about to move in or

out from the computer. The data size will affect the host I/O performance. The I/O process needs high bandwidth in order to have a stable throughput. The second task would be reading the data in bytes. This kind of task will consume I/O because it needs to perform operation in the logical disk itself and the lower the number of disk operation, the higher the input and output per second capacity.

The memory intensive task usually needs to perform analysis or searching in large scale of data. Big data application is the new trend of memory intensive task. This task usually needs to perform data mining and data analytic that involve a large dataset. During the process of data analytics, it typically starts with some large raw datasets, and then the applications will transform, clean and prepare the dataset for modelling with some sort of SQL-Like transformations. This will eventually utilize all memory that is assigned in the host.

In conclusion, different type of tasks can affect the performance of the virtualization. There is a need to conduct a performance analysis in virtualization that focuses on compute, memory and I/O, to improve the performance of these virtualized systems.

1.2 Problem Statement

Almost all of the scientific workflows have increased the number of Cloud computing paradigm, the techniques for virtualization such as low-overhead are becoming indispensable (Plauth, 2017). Although the container virtualization is a lightweight approach to execute these applications, it is very challenging in this unique environment to comprehend what will affect the workflow performance as it requires a lot of information and knowledge of the ever-changing application structure at any given moment. Based on (Cedia, 2017), the hardware resources of Cloud computing are always

limited, for this reason it is important that the available resources are adequately allocated according to the application behaviour to obtain the best possible performance.

Furthermore, such Cloud platforms have for most part were deployed to web-based platform and production applications, there is some portion that has dependencies and need to be overseen to run scientific workflows, particularly all the data intensive scientific workflows on the Cloud. Cloud computing provides an outlook changing and computing standard in terms of the extraordinary size of datacentre-level resource pool and provisions the on-demand resource mechanism, capable of addressing a large scale of scientific problems to capacitate scientific workflow explications.

Virtualization environments are flexible and unique by design, rapid changes on regular. Deploying a scientific workflow is risky, as minor differences in library versions on different servers can break the functionality of the workflow.

According to (Felter, 2014), every virtual machine that runs on Linux operating system have their own process and resource management abilities. One of the resource management is scheduling that is exported to the virtual machines. This reduce the administration and execution time but complicates resource management within the guest operating system. Every scientific workflow has its own requirements for it to run optimally. A new layer of complexity is added and can cause unfamiliar problems. Furthermore, each scientific workflow will consume the compute, memory, I/O and network of its host.

1.3 Objectives

This study aims to propose an automated system architecture that was able to choose the best virtualization platform to execute the workflow tasks. The benchmark performance evaluation was conducted on various workflow tasks, running on hypervisorbased and container-based virtualization to improve the selection decision.

The objectives of this research are outlined as follows:

- To conduct performance evaluation of workflow tasks running on container-based and hypervisor-based virtualization.

- To design an automated system architecture that choose the suitable virtualization platform for different workflow tasks.

- To evaluate the performance of various type of scientific workflows running on those virtualization platforms.

1.4 Scope

The research only focuses on hypervisor-based and container-based virtualization. The workflow tasks used for the evaluation are mainly compute-intensive, memory-intensive and I/O intensive. The tools used for the evaluation consist of open source and commercial tools. These tools must be able to measure the metric required and return the output of the performance test.

1.5 Dissertation Organisation

This section provides a general overview of the chapters that make up this dissertation. Chapter 1 contains the introduction to the topic in which the research is concerned with, and the problem statements are outlined. The objectives of the study, the scope as well as the structure of the dissertation are outlined. Chapter 2 describes the literature relevant to the research topic. First, the hypervisorbased and container-based virtualizations are discussed. Then, the related works, focusing on the performance evaluation are reviewed. After that, the performance tools that used for the research are explained. Subsequently, the study of workflow orchestration, scheduling and development are briefly discussed. The chapter ends with a summary that concludes the findings.

Chapter 3 discusses the research techniques used in the dissertation. The chapter begins with an introduction of the research flow that was used to conduct the research. The research phases, such as information gathering and analysis, proposed method, system design and implementation, system evaluation and documentation are discussed in detail in the subsequent sections of the chapter.

Chapter 4 discusses the system design, implementation and testing of the proposed technique. This chapter begins with the discussion of the proposed system architecture. Then, the system implementation and testing are presented.

Chapter 5 explains the process of evaluating the proposed techniques and a detailed discussion of the results. The performance tests were carried out in two different environments, namely hypervisor-based and container-based virtualizations. The chapter ends with a summary.

Chapter 6 concludes the research findings and achievements. The chapter begins with a discussion on how the objectives were obtained. Then, the chapter presents the research significance. Subsequently, the limitations and suggestions for future work are discussed.

CHAPTER 2: LITERATURE REVIEW

This chapter starts with a discussion about the virtualizations. There are two types of virtualization technology, namely hypervisor-based and container-based. Then, the chapter continues with a discussion on performance evaluation techniques and the performance tools. After that, the resource management concept is discussed. Subsequently the workflow is deeply briefed, and the types of workflow are discussed. Lastly, the chapter ends with a summary.

2.1 Virtualizations

Kumar (2015), said that virtualization plays an important role in Cloud computing and becomes the important key to the cloud infrastructure, as the technology can be enabled, the underlying hardware and software that are complex can be created as an intelligent abstraction layer hidden in the environment itself. There are two types of virtualization technologies namely the hypervisor-based and container-based virtualization technology. In general, there are several kinds of use cases for virtualization and basically both container-based and hypervisor-based virtualizations technology have strengths and weaknesses based on the respective use cases.

Anish Babu (2014) mentioned, there are three kind of virtualization technology which is Para virtualization, Container virtualization and Full virtualization. Para virtualization is a technique in which the guest operating system is aware that they are operating directly on the hypervisor instead of the underlying hardware. In Para virtualization, the virtualization supporting hypervisor is installed on the host operating system which runs over the underlying hardware. The second one is the Container virtualization. The Container virtualization is a technique in which each operating system kernel is modified to load multiple guest operating systems. Here guest operating systems are packed in the form of containers and each container will be allowed to load one by one. The kernel provides proper management of the underlying resources to isolate one container activity from the other. This type of virtualization technique has less overhead in loading the guest operating system in the form of containers and each container has their own IP address, memory and root access. The last one is Full virtualization. In Full virtualization, hypervisor supporting the full virtualization technique is installed directly over the underlying hardware. This hypervisor is responsible for loading the guest operating systems. And each guest operating system will run as if they are operating directly on the underlying hardware. That is each guest operating system will get all the features of the underlying hardware. Here hypervisor will directly interact with the underlying memory and disk space and hypervisor will isolate the activities of one guest operating system from the other. Hypervisors supporting full virtualization have a virtual machine management console from which each guest virtual machines can be easily managed.

2.1.1 Hypervisor

Desai (2013), said that Hypervisor allows user to run multiple OS on the same hardware and it will abstract the software layer from the OS. The virtual machine needs to be run on a hypervisor, and the computer that runs the hypervisor is called a host machine, and each of the VM that runs inside the host machine is called a guest machine. According to Merkel (2014), Hypervisor manages physical computing resources and makes isolated slices of hardware available for creating VMs. The VM creation are possible in the hypervisor environment as it split out the hardware resources into slices and the hypervisor will manage all the physical computing resource. The hypervisor can allocate resources to the respective virtual machine on demand. Hypervisor can be installed in two ways, the first one is installed directly on the hard disk of the computer and boot directly. Merkel (2014), also said that the other one is the hypervisor will be installed on top of the host operating system. As an example, the hypervisor that installed on bare-metal is VMware ESXi and the hosted hypervisor is called the VMware Workstation.

According to Eder (2016), The reliable grouping of a complete OS will be provided by the hypervisor-based virtualization, as the container itself, it is more focus on separating the processes from other process parallelly reducing the resource costs.

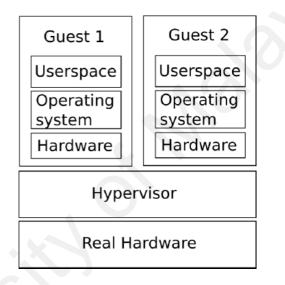


Figure 2.1. Scheme of hypervisor-based virtualization (Michael Eder, 2016)

Figure 2.1 shows that the scheme of hypervisor-based virtualization. Hypervisor can be installed directly to the hardware, and then on top of that the virtual machine can be deployed. The guest represents the virtual machine that sits on the hypervisor platform. Every VM has its own hardware settings and standalone OS running. The middle layer is where the hypervisor is installed. The hypervisor controls and manages all the running virtual machine. Hypervisor-based virtualization allows the imitation of another PC and copy different sorts of gadgets (for instance a cell phone), other computer models as well as other operating systems. This technology also takes advantage of modern compute capabilities. Besides, it will allow the application to directly access the compute, and virtual machine would receive the same benefit as in unprivileged mode. Therefore, it will result in the increase of the performance without sacrificing the host system security. After the provision of the virtual machine is completed, the application can be installed into the provisioned virtual machine.

2.1.2 Containers

Docker container is an open source software development platform. The main advantage of docker is the application can be grouped into a "container", enable it to be movable to any of the virtualization or bare-metal platform that runs Linux OS. A container is made from multiple small and non-heavy images, and each image is a template and convertible file system that includes all the middleware, libraries and binaries to execute the application. Pahl (2014), claimed that in the case of more than one images, the file system that supports read-only are stacked on top of each other to equally distribute the writable higher-level file system. The use of Docker containers will optimize existing apps while accelerating the way of applications delivery. With the hybrid portability, container will eventually eliminate the frictions of building migration plans from one source or platform to another. Moving container is seamlessly to the new cloud or new server as the container will separate column with their dependencies. In terms of security, the container can be applied to traditional application to decrease the attack from the surface layer, and at the same time mitigating risk and continuously monitor for vulnerabilities.

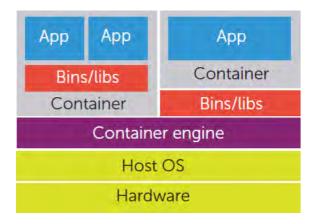


Figure 2.2. Container-based architecture in terms of managing its operating system (C. Pahl, 2014)

Figure 2.2 shows the container-based architecture. A container hold packages and individually-contained, custom pre-deployed parts of applications, and the business logic, such as binaries and libraries and the middleware. In a virtualized infrastructure, the sharing of the hidden platform and environment must be provided in a safer way. Containers can meet these necessities; hence, a more inside and out elicitation of specific concerns is required.

2.2 Performance Evaluation

Antonio (2011), mentioned that the main factor in research for computer architecture is performance evaluation. As the complexity grows, the complexity rate of the tasks that being executed by the computer system also rapidly increases for each task and programs. So, in order to measure the performance of the virtualization platform, speed is the best metric for developing an effective performance evaluation technique. The evaluation process will have a dependency based on its optimal trade-off. According to Pahl (2015), the main figure that runs the infrastructure layer is the virtual machines as it provides a virtualized OS. Container concept is similar as virtual machine but as it is using a lightweight technology, container will consume less resources and minimise the execution time. Low-level construct is provided by both virtual machine and container. Bernstein (2016), mentioned that the developer can navigate the operating system with an interface as interaction between the developer and the operating system. However, if the developer wants to deploy application in the Cloud infrastructure, it is recommended to deploy using container virtualization technology as it will package the application into the lightweight container.

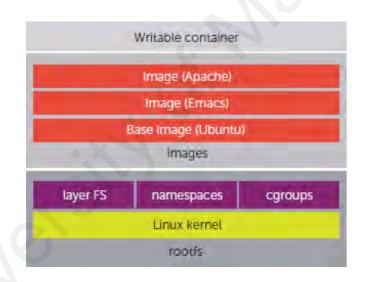


Figure 2.3. Container image architecture based on namespace and cgroup extension (Pahl, 2015)

As shown in Figure 2.3, the new Linux dispersions give part components, such as namespaces and control groups, to separate procedures on a combined operating system, support through the Linux Container (LXC) project. By separating the namespace will allow the groups of procedures to be isolated, keeping them from seeing assets in different

groups. Container innovations utilize distinctive namespaces for process detachment, network interfaces, access to network, mount focuses, and for isolating kernel and identify versions. Control groups oversee and constrain resource from getting through in process groups through extreme authorisation, accounting and, for instance, by restricting the memory accessible to a specific container as per said by Amit (2016).

Wang et al., (2017), mentioned that the main challenges of executing critical application within Cloud environment is the execution must satisfy the deadlines and response time in virtualization environment. This factor needs to be considered to secure guaranteed performance of the infrastructure. Therefore, performance evaluation on the applications is essential to obtain the most effective infrastructure platform for application deployment.

Based on Wang et al., (2017), different performance levels from the deployment of task will lead to different virtual machine services, as a result, it will obtain a different impact on the application cost and quality of services. Different requirements from the individual application will diverse the quality of services of the Cloud applications. Although the accuracy of the application is guaranteed, the application failure will lead to violation of the deadlines. Hence, execution timing for critical application should be carefully planned and maximise in the Cloud infrastructure. There are many perspectives that can be used to compare the performance of virtualization mentioned by Kozhirbayev (2017).

According to Zhanibek (2016), to evaluate container-based technologies from the perspectives of their overhead, it is crucial to measure the overheads incurred based on non-virtualized environments. The analysis conducted was focused on a range of performance criteria: the performance of compute, memory, network bandwidth, latency and storage overheads. In order to obtain accuracy and consistency in the bunch of results,

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a lot of experiments were repeated, in which the average timing and standard deviation were recorded.

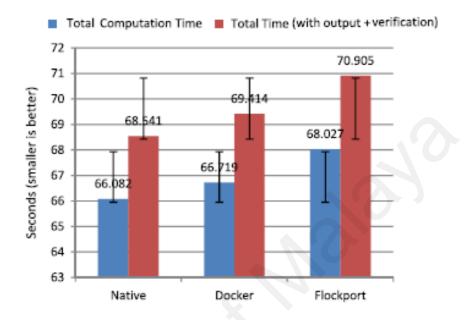


Figure 2.4. Performance result of the Y-cruncher benchmarking tools (Kozhirbayev, 2017)

Figure 2.4 shows that Y-cruncher, one of the benchmarkingtools that was used to test the compute tasks and generally used, as a test for multi-threading tools running in multicore systems. Y-cruncher can also be used to calculate the Pi value and measure the gap of other constants. The metric such as the total time executed, computational time and multi-core efficiency can be variously performed by the Y-cruncher. From the figure, bare-metal or native environment perform similarly as the Docker based on the computational time, as for the Flockport, it took 1.3 seconds longer.

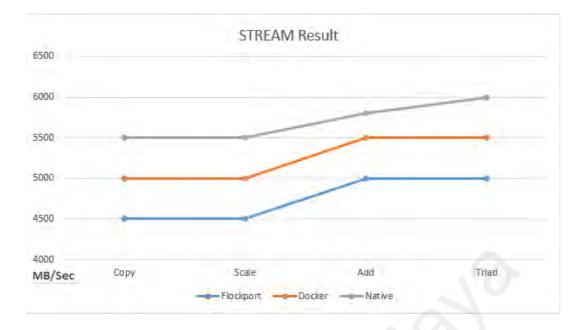


Figure 2.5. STREAM result (Kozhirbayev, 2017)

As for the memory performance evaluation, the STREAM software is being used to test the micro-hosting environment. STREAM determines the throughput of the memory that being utilized straightforward from vector kernel procedures. Figure 2.5 shows that Docker produces marginally better result than Flockport and not much different with the native platform.

The study from Helen Karatza (2017), measured the overhead caused by virtual machines while containers are running on top of them. A series of benchmarks were tested to measure the additional layer of the VM affects the CPU, memory and network performance. In order to allow them to use all available resources, the measurement of CPU overhead were contacted with LINPACK.

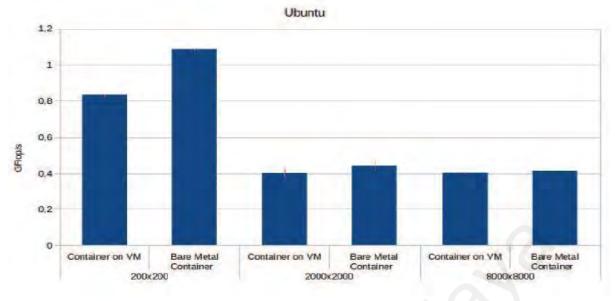


Figure 2.6. LINPACK on Ubuntu (Helen Karatza, 2017)

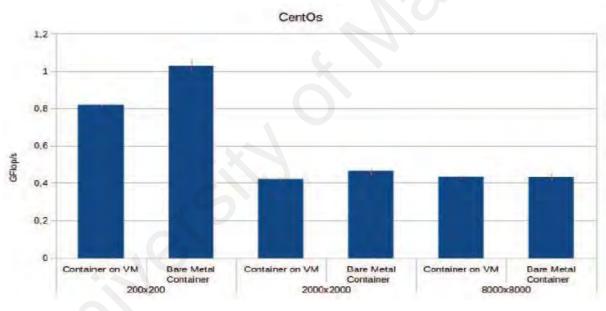


Figure 2.7. LINPACK on CentOS (Helen Karatza, 2017)

Figure 2.6 and 2.7 shows that the LINPACK result tested on Ubuntu and CentOS. The LINPACK was run in containers on bare metal and in containers on top of a VM for matrix sizes 200x200, 2000x2000 and 8000x8000. As per observations, LINPACK measured the best performance with the smallest table (200x200). In the case of the smallest table, we noticed the biggest overhead which was caused by the additional

virtualization layer of the VM and is about 28.41% in average for the two operating systems. This happens because VMs hide the nature of system information from the execution environment. Regarding the bigger tables, we observed a lower performance, as well as a much smaller gap (0.87 % in average for the two operating system) between the container on bare metal and the container on top of the VM. This probably happens because the necessary data cannot be cached in the available high speed memory which as a result increases the overhead.

The benchmark tool that were used to evaluate the I/O performance is STREAM. It executes four different operations namely Copy, Scale, Add and Triad. This benchmark intends to measure the main memory bandwidth and not the cache bandwidth, however STREAM recognizes a strong relation between the evaluated throughput and the size of the CPU cache.

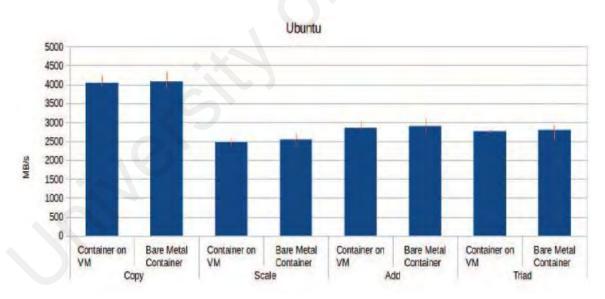


Figure 2.8 STREAM on Ubuntu (Helen Karatza, 2017)

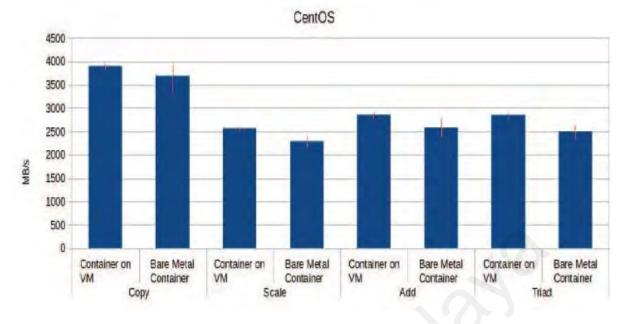


Figure 2.9 STREAM on CentOS (Helen Karatza, 2017)

In Figure 2.8 and 2.9 shows the STREAM measurements on four different operations. We generally observed that the overhead caused by the additional virtualization layer of the VM is about 4.46% in average regarding the Copy operation, 1.76% in average regarding the Scale operation, 2.39% in In Random Read and Random Write, data accesses are carried through in random locations within the file and are affected by factors such as the OSs cache and seek latencies. With Mixed Workload we measured the performance of reading and writing operations combined and we found a throughput reduction by 14.67% affected by the VM layer.

The Network I/O performance measured with NETPERF benchmark. NETPERF can take unidirectional throughput and end-to-end latency measurements. In order to run the experiments we used another physical node to host the NETPERF server, whereas the client ran on our initial physical node. The TCP STREAM is used to measures the throughput of transmitting TCP packets of 4, 16, 64, 256 and the default 16384 bytes between the NETPERF client and the NETPERF server. We also took TCP RR and UDP RR measurements. TCP RR is a process of multiple TCP requests and responses in the

same TCP connection that is common in database applications and UDP RR that uses UDP request and response packets.

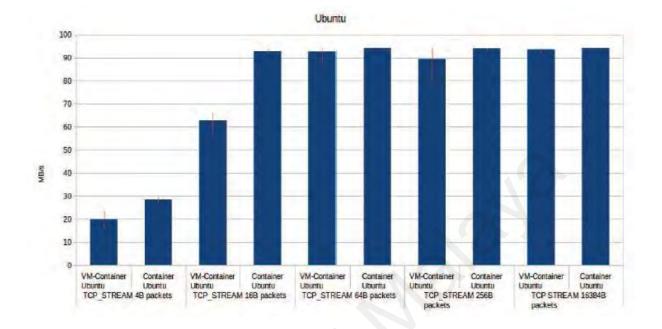


Figure 2.10. NETPERF, TCP_STREAM on Ubuntu (Helen Karatza, 2017)

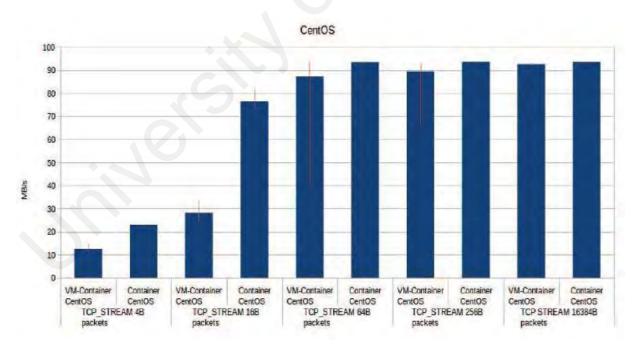


Figure 2.11. NETPERF, TCP_STREAM on Ubuntu (Helen Karatza, 2017)

Figure 2.10 and 2.11 present the TCP STREAM measurements for the default and 5 different packet sizes for the. The two operating systems running on bare metal and on top of a VM. In all cases we observe that, regarding the smaller packet sizes, the overhead that the additional layer of VM virtualization causes to the network is significantly higher compared to the larger packets' overhead. This is quite reasonable because of the smaller packets require more computation power. For the small 4 bytes packets there is a mean throughput reduction by 33.3% in both operating systems, whereas, regarding the bigger size of 16384 bytes (default packet size) there is a mean throughput reduction by only 0.69% in both operating system when containers run on top of VM.

2.2.1 Performance Evaluation Tools

There are some tools that are used to measure the performance effectiveness between virtual machines and Docker containers. These tools were specifically designed to perform different kinds of test, such as compute, memory and I/O.

2.2.1.1 Bonnie++

Bonnie++ is a performance tool that allows the benchmarking of how the filesystems perform numerous tasks, that makes it a valuable tool once changes are made to the RAID, how the filesystems are created, and how the network filesystems perform. Bonnie++ benchmarks three things, which are data searching and write speed, range of seeks that may be performed per second, and range of file metadata operations that are able to be performed per second. Metadata operations contain file creation and deletion as well as obtaining metadata, like the file size or owner. As the test is performed, Bonnie++ prints a line, informing that the test is executed. The Bonnie++ distribution also includes the

html page to display all results. It will generate a table for the output after performing the test.

2.2.1.2 SysBench

SysBench is a standard, cross-platform and multi-threaded benchmark tool for evaluating OS parameters that are vital for a system running an information beneath intensive load. SysBench is a benchmark suite which gives a quick impression about the system performance. The idea of this benchmark suite is to quickly get an impact concerning about system performance while not setting up advance database benchmarks. The design is extremely easy. SysBench runs a given range of threads, within which all execute requests in parallel. The workload created by requests depends on the required pilot mode. Either the whole range of requests or total time for the benchmark are often restricted, or both. Obtainable test modes are implemented by compiled-in modules, and SysBench was designed to simplify the addition of new test modes. Each test mode might have additional (or workload-specific) choices.

Table 2.2 shows the types of performance evaluation test that can be performed and the output that is obtained from the test. As stated, there are four types of performance evaluation test and each of the test will collect different kinds of output and result.

2.2.1.3 Y-Cruncher

Y-cruncher is a program that can compute Pi and other constants to trillions of digits. It is the first of its kind that is multi-threaded and scalable to multi-core systems. Ever since its launch in 2009, it has become a common benchmarking and stress-testing application for overclock and hardware enthusiasts.

The main computational features of y-cruncher are:

- Able to compute Pi and other constants to trillions of digits.
- Two algorithms are available for most constants. One for computation and one for verification.
- **Multi-Threaded** Multi-threading can be used to fully utilize modern multicore processors without significantly increasing memory usage.
- Vectorized Able to fully utilize the SIMD capabilities for most processors.
- **Swap Space** management for large computations that require more memory than there is available.
- Multi-Hard Drive Multiple hard drives can be used for faster disk swapping.
- Semi-Fault Tolerant Able to detect and correct for minor errors that may be caused by hardware instability or software bugs.

2.2.1.4 STREAM Benchmarking Tools

The STREAM benchmark is a simple synthetic benchmark program that measures sustainable memory bandwidth (in MB/s) and the corresponding computation rate for simple vector kernels. It is specifically designed to work with datasets much larger than the available cache on any given system, so that the results are more indicative of the performance of very large, vector style applications.

For this research, instead of using STREAM and Y-cruncher, Sysbench and Bonnie++ will be selected as the performance evaluation tools. This is because, for I/O testing, custom file size is allowed by Sysbench. Hence, evaluation testing will be more comprehensive and accurate. As for STREAM, during the evaluation test, the total number of threads cannot be defined initially. Furthermore, the STREAM output did not show the statistic of minimum, average and maximum output per-request. Sysbench on the other hand, the result shows all the mentioned with clear and informative way. Compared from Bonnie++, Y-cruncher is also multi-threaded program, but it is deterministic. It was designed to avoid problem that common occurs in asynchronous application and can be extremely difficult to track down and fix. Because of this determinism, testing crashes and errors happen intermittently. Bonnie++ is very flexible as the file size can be set from 1GB onwards. The output provided also were split into 2 portion which is the command output and the HTML table output. There is also information regarding how many I/O blocks transferred per second. That is the reason Bonniee++ and Sysbench is being used for this research.

Type of Performance Evaluation Test	Purpose
Compute	To examine how many units of information a system
•	processes over a specific time.
Memory	To evaluate the memory working storage space available
	to a processor or workload.
I/O	To observer the response time or latency based on the
	amount of time that elapses between every read and write
	job or task.
Network	To check the bandwidth or the volume of data per second
	that can move between workloads and usually across
	networks.

Table 2.2. Main purpose of performance evaluation test

2.3 Resource Management

Gangadhar (2018), claimed that resource management is an essential technique to utilize the underlying hardware of the Cloud efficiently. Resource management is

important in virtualized infrastructure as it gives a clear picture on the amount of task and job that must be done and helps to schedule and plan the task executions by allocating suitable resources for every task and job. Resource management provides an efficient and effective deployment of a job and task with a satisfying result. The purpose of resource management is to manage the scheduling task for all the jobs before they are executed. This portion involves allocating resources into the job.

According to Linthicum (2016), the goal of managing the application-tier and deploying application designs most likely can be achieved by containers. On the operating system level, container can manage the applications such as Web servers, database servers and application servers. Furthermore, the hypervisor-based platform is focusing to separate and hold the resources based on machine-by-machine. As for containers, the CPU resources can be shared and distributed than VMs.

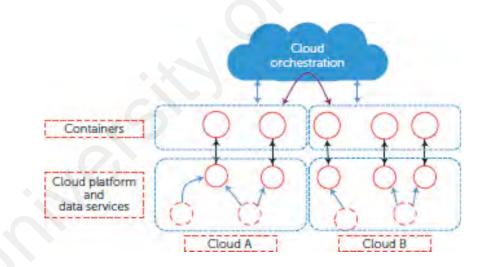


Figure 2.12. A Cloud orchestration layer oversees the infrastructure supporting live migration of containers (Linthicum, 2016)

Figure 2.6 shows that containers can be migrated or moved from Cloud to Cloud orchestration layer. All the containers can be run, leveraged and automatically migrated

from one Cloud to another, to support the infrastructure requirements. Computing capabilities can be distributed since the applications can be separated into various kind of domains.

Non-stop monitoring of resources is vital in managing the virtual environment mentioned by Pooja & Pandey (2014). To guarantee the Service Level Agreement (SLA) while optimally consuming the resources, performance evaluation needs to be considered in order to maximise the utilization of the compute system. As a result, different users can share the same single multicore node claimed by Miguel et al., (2013). The best virtualization that can increase the percentage of resource sharing is the container-based virtualization by allowing multiple separated instances of user-space. Meanwhile, the disadvantage of container-based virtualization is it cannot firmly isolate the resource as good as the hypervisor-based virtualization.

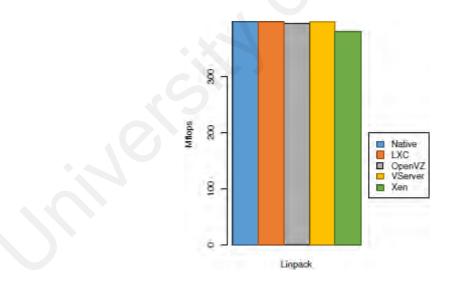


Figure 2.13. Computing performance by using LINPACK for matrices (Miguel et al., 2013)

Figure 2.7 shows that LINPACK was used as a benchmark to evaluate the computing performance on a single computer-node. Miguel et al., (2013), claimed that

LINPACK consists of a set of subroutines that analyses and solves linear equations by the least square's method. Result of the LINPACK can be used to estimate performance and it is run over a single processor. LINPACK is used as a tools to measure the matrices of order 300 in all container-based systems and compare them with Xen. The first one is Linux-VServer, this system is one of the oldest implementations of Linux container-based system. Instead of using namespaces, Linux-VServer introduced its own capabilities in the Linux kernel, such as process isolation, network isolation and CPU isolation. This system also does not virtualize network subsystem as all the networking subsystem within all containers share the same routing and tables IP. The second system is the OpenVZ. This system offers similar functionality to Linux-VServer. However, it is built on top of kernel namespaces, making sure that every container has its own isolated subset of a resource. Moreover, the OpenVZ uses the network namespace. In this way, each container has its own network stack, which includes network devices, routing tables, firewall and so on. The third system is the LXC. In the same way as OpenVZ, LXC uses kernel namespaces to provide resource isolation among all containers. Furthermore, unlike the OpenVZ and Linux-VServer, the LXC only allowed resource management via cgroups. Thus, LXC uses cgroups to define the configuration of network namespaces. As for Xen system, it will be use as the representative of hypervisor-based virtualization, because it is considered one of the most mature and efficient implementations of this kind of virtualization. As the result, there was not much difference obtained from the experiment as all the container-based systems obtained similar result to the native, as shown in Figure 2.7. It is due to the fact that there was no influence of the different compute schedulers when a single compute -intensive process is run on a single processor.

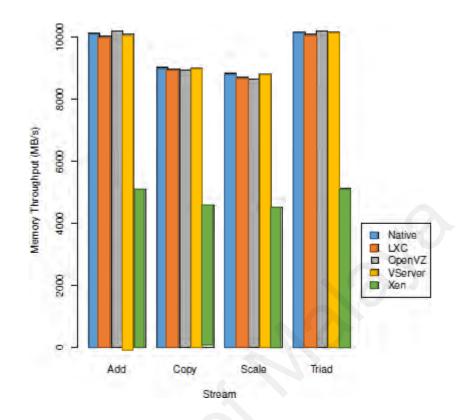


Figure 2.14. Memory throughput by using STREAM (Miguel et al., 2013)

As shown in Figure 2.8, the memory performance evaluation is evaluated with STREAM, a simple artificial benchmark program that measures the properties of memory bandwidth. This performance evaluation tool will perform four types of vector operations, which is copy, add, scale and triad, and uses a larger dataset than the cache memory available in the infrastructure. As observed, the result was similar as native system for container-based due to the unutilized memory of the host, enabling a better use of memory. Unfortunately, the worst result was in Xen, which presented an average overhead of approximately 31% as compared to the native throughput. The hypervisor-based virtualization layer implements the translation of memory accesses that causing the memory to be overheaded and resulting a decrease in terms of performance.

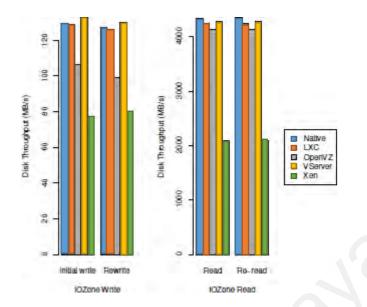


Figure 2.15. Disk throughput by using IOZone (Miguel et al., 2013)

Figure 2.9 shows the result of disk performance evaluation by using IOZone benchmark tool. The IOZone can generate, access pattern and measure a variety of file operations. A record size of 10GB and 4KB will be run as a test case. According to graph, it revealed that both Linux-VServer and LXC had a similar result for read and re-read operations. Furthermore, a gain of performance was achieved as compared to the OpenVZ results which is lower than both Linux-VServer and LXC. The root cause probably was due to the I/O scheduler was being used by the different systems. The worst result was obtained in Xen for all I/O operations as caused by the para-virtualized drivers. These drivers weren't ready to attain a high performance nevertheless.

Furthermore, in an observation-based study by Jeroen (2014), the basic principle of a container was that it allowed for processes and their resources to be isolated without hardware separation or hardware dependencies. Containers provide a form of virtualization platform that each container will run their own OS sharing the kernel.

From the above findings, it is very important to conduct a performance evaluation on different approaches of virtualization. The performance results can help both researchers

and practitioners to design better virtualization architecture according to application needs. The study can be viewed as a foundation for more sophisticated evaluation in the future.

2.4 Docker-Sec Automation Architecture

Docker-sec is an open-source, automated and user-friendly mechanism for securing Docker and generally OCI2 compatible containers. Docker-sec offers users the ability to automatically generate initial container profiles based on configuration parameters provided during container initialization. If a stricter security policy is required, Dockersec can dynamically enhance the initial profile with rules extracted through the monitoring of real-time container execution. Docker-sec adds an additional security layer on top of Docker's security defaults by automatically creating each container AppArmor profiles claimed by Fotis Loukidis (2018).

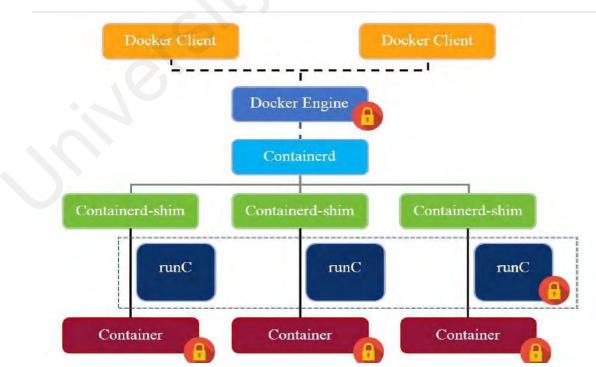


Figure 2.16 Docker Components Protected with AppArmor in Docker-Sec

(Fotis Loukidis, 2018)

Figure 2.10 Show how the architecture of Docker component with AppArmor. Docker-sec creates secure AppArmor profiles for all Docker components for rendering the environment. Container profiles are created automatically using rules extracted from the configuration of each container and enhanced with rules based on the behaviour of the contained application. To that end, Docker-sec employs two mechanisms which is static analysis, which creates initial profiles from static Docker execution parameters and dynamic monitoring which enhances them through monitoring the container workflow during a user-defined testing period. The goal is to construct a separate profile per container, placing each one in a separate security context in order to restrict the sharing of resources among containers. The components of Docker that are automatically protected via AppArmor profiles via Docker-sec are designated with red lock logos.

2.5 Workflow

Single or multiple computational task can be easily expressed by users using the scientific workflow. This computational tasks include reformatting the data, running an analysis from an instrument or a database. The dependencies of the task in most cases can be described by scientific workflow as a directed acyclic graph (DAG), where the edges denote the task dependencies and the nodes are tasks. Scientific workflow is very efficient is managing the data flow. Everything from very large parallel task to short serial task surrounded by small, serial tasks used for pre- and post-processing as mentioned by Pegasus (2013).

According to Paul Martin et al., (2016), a common strategy to make the simulation experiments more manageable is to model them as workflows and use a workflow management system to organise the execution. The automation process of the business in which task, documents or information are passed from one participant to another can be defined as a workflow, based on the set of its procedural rules said by O' Brien (2010).

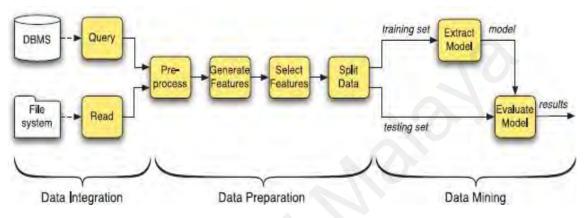


Figure 2.17. Common workflow in scientific experiments (Martin et al., 2016)

As shown in Figure 2.10, a workflow was split into three components. The first one would be a list of operations or tasks, the second is a group of dependencies between all the internal connected tasks and finally the group resources that contains data is used to execute or remove the flow. Basically, from the figure above, the vertices are the data and tasks resources. The edges dependent on the connection of the vertices. The edges can represent two kinds of dependency, which are the control flow and data flow. Control flow graphs involve task and priority requirements. The data-flow depends on the conditions of the undertakings flow of data. The data move along into a circular segment and are changed by the pre-process operator. Every data item will be changed by the pre-process operator and the result will be transmitted to succeeding operator. The graph of the data-flow allow the operators to execute to cover in a prepared pipeline.

2.5.1 Workflow Orchestration

This workflow orchestration will separate the tasks based on their intensiveness. The intensiveness falls into 3 categories, which is compute, memory and I/O intensive. This orchestration will arrange the task in sequence and manage the task before it is scheduled.

2.5.2 Workflow Scheduling

Workflow scheduling schedules the task to be deployed on the respective platform. This scheduling will be decided based on the workflow intensiveness and different category will be deployed in a different platform in order to receive the fastest deployment and response time.

2.5.3 Workflow Deployment

After all tasks were sorted in sequence and scheduled, the deployment is where the task will be executed into the platform. The task will be executed in hypervisor-based or container-based virtualization platform. Table 2.3 shows the literature summary.

Author	Year	Analysis	Summary
Claus Pahl	2014	Understanding the concept	Virtual machines (VMs) are the
		of virtual machines and	backbone of the infrastructure
		container in virtualization	layer, providing virtualized
		technologies are important	operating systems (OSs).
		to have a clear view of what	Containers are similar but have a
		is needed to be evaluate.	more lightweight virtualization

			concept with less resource and can
			minimise time-consumed.
David	2016	Study the suitable	Both VMs and containers provide
S. Linthicum		operating system to run on	a rather low-level construct
		both virtual machine and	Basically, both present an OS with
		container is crucial to have	a GUI to the developer.
		a better performance	
		comparison result.	0
Junchao	2017	Analyse the best metric to	Execute time crucial applications
Wang et al.		be use as a benchmark to	at intervals of Cloud environments
		compare both virtual	whereas it satisfies execution
		machine and container	deadlines and interval time
		performance.	necessities could be a challenge as
			it will be difficult in securing
			guaranteed performance from the
		S	underlying virtual infrastructure.
Miguel et al.	2013	To Study the core objective	Resource management core
		of resource management in	objective is to maximise the overa
		order to build a proper	utilization whereby multiple
		resource management	number of users can share a single
		architecture.	multicore-node
Fotis	2018	Study the disadvantages of	Containers pose significant securit
Loukidis		using containers to avoid	challenges due to their direct
		any irrelevant experiment	communication with the host
		test.	kernel, allowing attackers to break
			into the host system and locate

			containers more easily than virtual
			machines.
Helen	2017	Survey the feature	VMs and containers have different
Karatza		differentiation on both	features. The main advantage of
		hypervisor and container	Container is low performance
		virtualization technology to	overhead whereas VMs
		understand which platform	present strong isolation. To
		would be the best for	enhance security, containers may
		running specific workflow	need run on top of a virtual
		applications.	machine.
Anish Babu	2014	Study the type of	There are mainly three kind of
		virtualization technologies	virtualization technologies being
		to understand the different	used which is Para virtualization,
		and choose the right	Container Virtualization and Full
		platform to compare with	Virtualization. Each virtualization
		container technologies.	techniques have different
	0		implementation
		V	

2.6 Chapter Summary

This chapter discussed the research background of virtualization and resource management. Several applications and research done in this field were studied and explained in detail. Both hypervisor-based and container-based have pros and cons. The problem of the existing hypervisor and container virtualization technology were identified.

Table 2.4 shows the problems that were encountered when the workflow tasks were executed in the virtualization platform. Most of the impacts caused a decrease in performance and inefficient load balancing in the deployment of scientific workflows.

Table 2.4. Literature Summary

Problem Identified
Workflow Performance Decreases
Worknow I error mance Decreases
This performance impact happens mainly since there are multiple
memory and compute managers for a guest OS, causing an overhead and result
in a higher performance impact.
· ×
Poor Resource Allocation
This is because all workflow requires specific resource allocation for
execution. Therefore, a proper load balancer is needed to achieve the optimum
performance of the executed workflow.
Complexity of the Workflow
Scientific workflow is one of the complex workflows and the complexity of
this workflow will consume different compute, memory and I/O utilization of
its host if it does not deploy in a suitable platform.

To address the problems, an automate system architecture technique was proposed. The methodology adopted for this research work will be discussed in the next chapter.

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CHAPTER 3: METHODOLOGY

This chapter discusses the research techniques used in the dissertation. The chapter starts with a discussion about the research methods then, the research phases, such as information gathering and analysis, proposed method, system design and implementation, system evaluation and documentation are discussed. A summary is provided at the end of chapter.

3.1 DoKnowMe

DoKnowMe is a software engineering methodology that can be used to evaluate performance and has become one of the main factors in software quality assurance. According to (Li & O'Brien, 2016), DoKnowMe was employed to guide the study evaluation implementations. DoKnowMe is a part of evaluation methodology on the analogy of "class" in object-oriented programming. Driven by the concept of objectoriented programming, the general performance evaluation logic was distinguished, and an abstract evaluation methodology was developed into the term "Domain Knowledgedriven Methodology (DoKnowMe)". As the predefined domain-specific knowledge were remained, DoKnowMe can be summarised into more specified methodologies to help evaluate the computing system and different software performance. There are four factor that is a generic validation which is repeatability, usefulness, feasibility and effectiveness. All factors were used to validate the methodology in the evaluation domain. With good and promising evaluation results, more evaluation strategies can be integrated to improve DoKnowMe, and hence become more specific on the performance evaluation of virtualization technologies.

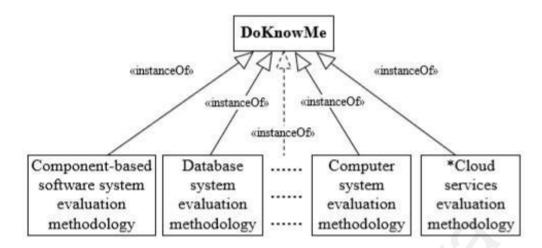


Figure 3.1. The relationship between DoKnowMe and its instance methodologies

Figure 3.1 shows the DoKnowMe methodology that can be used in the experimental measurement scenario. As a fundamental evaluation scheme, experimental performance measurement was considered to determine the annotations of performance involved in the model (Koziolek, 2010). The study of domain knowledge methodology is crucial. The evaluation of knowledge can be learnt from major experts and various publications.

The performance evaluation procedure consists of a sequential-process as well as recursive experimental activities. After the evaluation implementation is prepared, the recursive experimental activities will executed in a group of experimental tests. Then, the iteration of the experimental design is decided by the experimental results and analysis from the prior iteration.

Figure 3.2 explains the step-by-step procedure by using DoKnowMe. Every evaluation step is considered as the I/O component. DoKnowMe basically integrates and facilitates activities to essentially comprise the evaluation activities. Therefore, only some theoretical examples are executed to demonstrate particular evaluation steps that are involved in DoKnowMe.

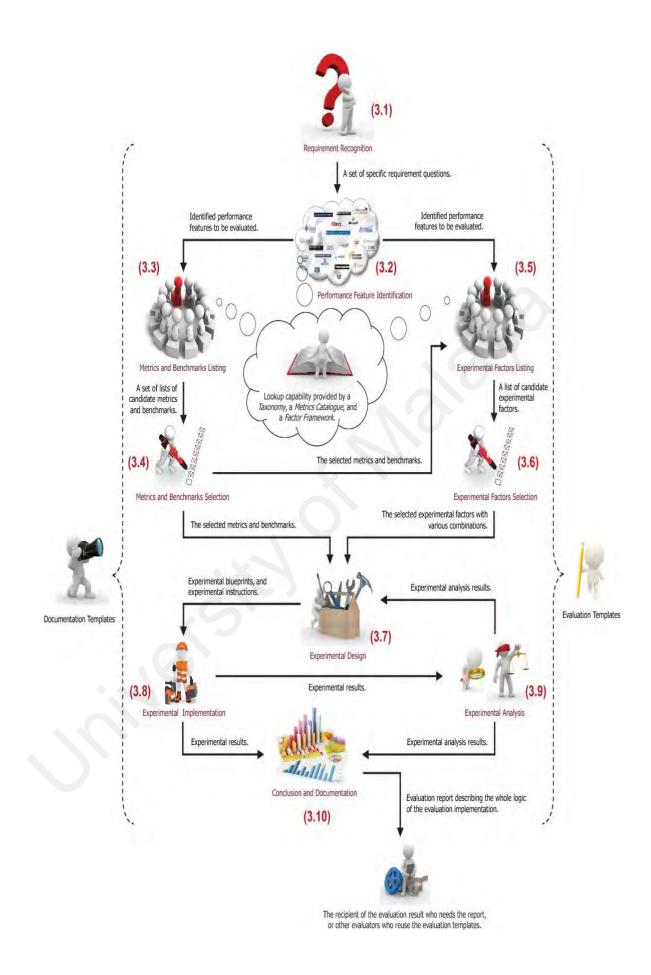


Figure 3.2. The step-by-step procedure by using DoKnowMe

Figure 3.2 shows the sequential procedure of DoKnowMe. All procedures involved in the methodology are considered on "what", "how", "why", "when" and "where" the experiment tests should occur. Each procedure has its own dependencies and all procedures must be in sequence so that the objective of the experiment can be achieved.

3.1.1 Requirement Recognition

The recognition of a requirement is not only to understand a problem related to the system performance evaluation, but also to realize a transparent statement of the analysis purpose, which must be clear with a non-trivial task. A clearly nominative evaluation requirement will facilitate correct driving of the remaining steps within the evaluation implementation. To solve performance evaluation, the evaluation requirement will be identified based on the research problem statements. The problem statements will find an idea on how to measure the performance effectiveness in both virtualization platforms (Zheng, 2016).

3.1.2 Performance Feature Identification

Given the clarified evaluation requirement of a system, evaluators need to further identify relevant performance features to be evaluated. Since different end users could be running different kind of application in the virtualization platform, it would be difficult for evaluators to directly locate proper performance features. Therefore, it will be helpful and valuable to measure the performance based on applications that runs on hypervisorbased or container-based virtualization platform.

3.1.3 Metrics and Benchmarks Listing

The choice of the right metrics depends on the identified performance features to be evaluated, while given that the performance features are yet insufficient for choosing right metrics. On the one hand, a performance feature can be measured by different metrics with different benchmarks. Therefore the performance metric that can be used to evaluate the performance is CPU, network, I/O and memory.

3.1.4 Metrics and Benchmarks Selection

A metric is a measurable quantity that precisely captures some characteristics of a performance feature. The selection of metrics plays an essential role in evaluation implementations. Hence, the metric that really effective to measure the performance of the virtualization platform is CPU, memory and I/O. So this three metric will be the benchmark of the performance evaluation.

3.1.5 Experimental Factor Listings

Although listing a whole scope of experimental factors may not be simply achieved, the factor listings should be kept as comprehensive as possible continually. This is often for more analysis and decision-making regarding the factor choices and information collected. Therefore, the aim of experimental factors listing is to list all experimental factors associated with performance features that are to be evaluated. In the present study, the compute, memory and I/O are used as the experimental metric. Hence, there are set of performance evaluation tools that we can chose to run the performance evaluation experiment such as Sysbench, STREAM, Y-cruncher and Bonnie++.

3.1.6 Experimental Factors Selection

There is no conflict between selecting limited factors in this step and keeping a comprehensive factor list in the previous step. On the one hand, an evaluation requirement usually comprises a set of experiments, and different experiments might have to select different factors from the same factor list. Hence, Bonnie++ and Sysbench is being chose as the performance evaluation tools to measure the CPU, memory and I/O intensiveness within the hypervisor-based and container-based virtualization technology.

3.1.7 Experimental Design

Once the experimental factors are selected, the evaluation experiments are ready and designed. Normally, a low scale of pilot experiment would profit the relevant experimental design, by serving evaluators to urge conversant in the experimental environment and optimize the experimental sequence. The experiment design involves the way the experiment should be conducted. The experiment will be performed in a workflow design. So basically, there are 4 sets of workflows that will be use in the experiment test. The first one is the CPU intensive workflow which the workflow will consist a lot of task that requires CPU. The second set would be the memory intensive workflow. The memory intensive workflow will also have a lot of task that requires memory from its host. The third set is the I/O intensive workflow. This set of workflow will have a lot of input and output task will requires its host to allocate more resources. The last set is the uniform workflow. This workflow will consist of CPU, memory and I/O task.

3.1.8 Experimental Implementation

Implementing an experiment involves carrying out a series of experimental actions, starting from the experimental environment preparation to running benchmarks. Generally, experimental results might answer their corresponding evaluation requirement queries. However, in most cases, a lot of convincing conclusions are drawn through further experimental analysis. So for the performance evaluation test, there a 4 sets of workflow experiment is conducted. All results are recorded and analysed during the implementation of the test cases to see if there is any difference in response time and application responsiveness in both hypervisor and container-based virtualization platforms.

3.1.9 Experimental Analysis

Experimental results can sometimes answer their corresponding evaluation requirement questions already. However, in most cases, more convincing conclusions would have to be drawn through further experimental analysis. So based from the test that were conducted, a statistic diagram is formed to ease to draw the conclusion and potential decision-making process.

3.1.10 Conclusion and Documentation

DoKnowMe uses a structured manner to implement conclusion and documentation. Tables and visual representations of experimental analysis results are used to interact to the pre-specified requirement queries. The answers to all requirement queries are summarised into linguistic communication findings to mirror the conclusions higher. Based on the results obtained from the experiment test, we can have a better conclusion from the report.

3.2 Chapter Summary

The DoKnowMe research methodology plays some key roles in evaluating performance of the virtualization technology. An automated system architecture is proposed in this research to assess and measure both virtualization platform and deploy application in the suitable platform. The next chapter elaborates the system design, architecture, implementation and testing phases.

CHAPTER 4: SYSTEM DESIGN, IMPLEMENTATION AND TESTING

This chapter discusses the system design, implementation and testing of the proposed technique. First, the proposed system architecture is discussed, followed by a discussion about the system implementation and testing. A summary is presented at the end of the chapter.

4.1 Automated System Architecture

This research work proposes an automated system architecture that could dynamically allocate workflows jobs to different virtualization platforms according to specifications. Figure 4.1 shows the proposed automated system architecture which consists of three layers. At the top layer, the workflows can be submitted with specifications and needs.

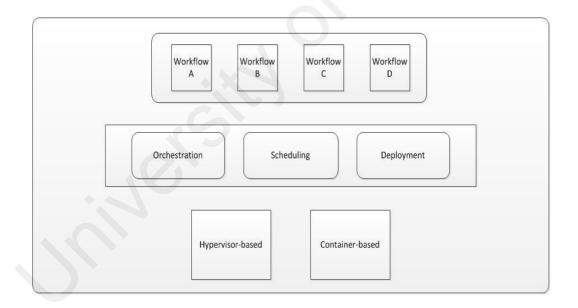


Figure 4.1. Automated system architecture

The middle layer is the core layer of the system architecture. It consists of three main components, namely orchestration, scheduling and deployment. The orchestration component is used to orchestrate the workflow process and define the workflow specifications. This component defines all jobs in the workflow together with their relations. Besides, each job parameters are stated in addition to the specifications. This information is essential for scheduling components. Scheduling is the second core component in the middle layer. Scheduling plays a vital role in terms of allocating the resources to the jobs in each workflow. The objective of scheduling function is to minimise the total execution time for each workflow.

The last core component is deployment. Deployment handles all workflows before they are deployed to the platform. In the hypervisor-based platform, deployment provisions the virtual machine and configures the virtual machine specifications based on workflow criteria. Meanwhile in container-based platform, deployment handles all the container creations and install all dependencies for the workflow application to run in those containers. Each container will extract library and bin files and provision an image for it.

The bottom layer consists of two types of virtualization platforms, namely hypervisor and container. The system will deploy the workflow application to the selected virtualization platform and return the output once the execution is completed. The virtualization platform is set up automatically based on workflow needs.

Figure 4.2 shows the interaction of the core components in the proposed system architecture. The applications are categorised into 3 categories, namely compute intensive, memory intensive, and I/O intensive. The workflow needs can be specified based on the above categories.

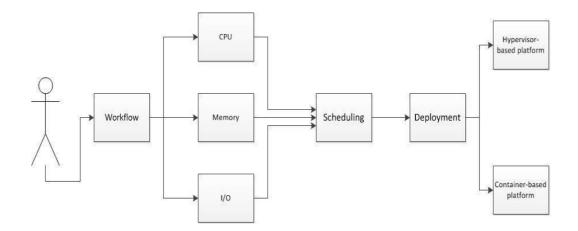


Figure 4.2: Workflow categorisations

Figure 4.3 shows the use case diagram of the proposed system architecture. The use case diagram provides a better understanding of the functionality of the system and the system interactions with users. As shown in the Figure 4.3, the workflow specifications can be specified during submission. Besides, results can be automatically collected from the system. The system will automate the platform selection process based on workflow specifications.

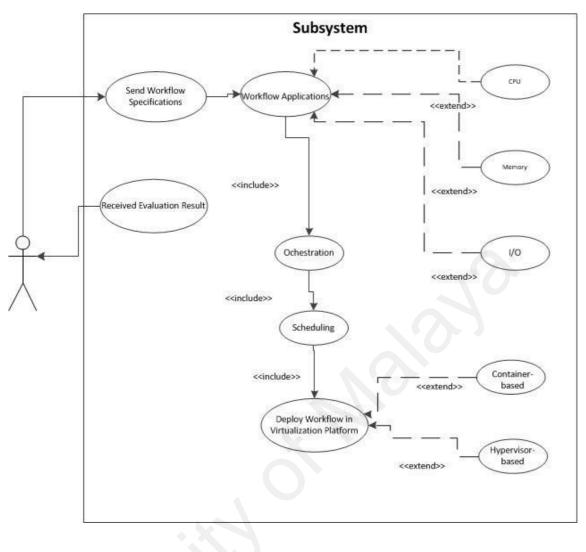


Figure 4.3. Use case diagram

Table 4.1 shows the detailed descriptions of the above Use Case Diagram. This table describes each of the case in the diagram. All the cases have their dependencies and requirements.

Table 4.1.	Use	case	descriptions

Use Case	Description
Send the workflow specifications	This use case is initiated by the end user to submit input of workflow specifications.
Receive the evaluation result	This use case is initiated by the automation system and the output will be forwarded to user so that the user will obtain the result.

	 Main flow of events: 1. The user sends the request and waits for the automation system decision 2. The system will evaluate the result and map it to the resource
Deploy workflow applications in suitable platform	This use case is when the workflow is already scheduled and orchestrated.
	 Main flow of events: 1. System will receive scheduling result in sequence 2. System will deploy workflow based on the scheduler decision.
Workflow Applications	This use case is initiated by the automated system after the input was received from the user. Main flow of events: 1. The system will wait for the input from the
	 user. 2. Automation system will orchestrate the workflow specifications based on user input 3. The system will schedule the workflow based on their priority. 4. After scheduling task completed, system will deploy workflow in the suitable platform. 5. System will provide output of the evaluation result to the user.
Virtualization Platform	This use case is initiated by the user based on the output provided by the automation system.
	 Main flow of events: 1. User will wait until the evaluation result is received by the automation system. 2. System will deploy the workflow based on the output provided by the automation system.

Figure 4.4 shows the system flow of the proposed system architecture. Each user is required to input the workflow name, workflow specifications in XML file and upload the executable files for each job in the workflow. Once the workflow is submitted, the system will orchestrate the workflow into different categories. Based on the information, the automate system architecture will schedule and deploy each job in the workflow to the suitable virtualization platforms.

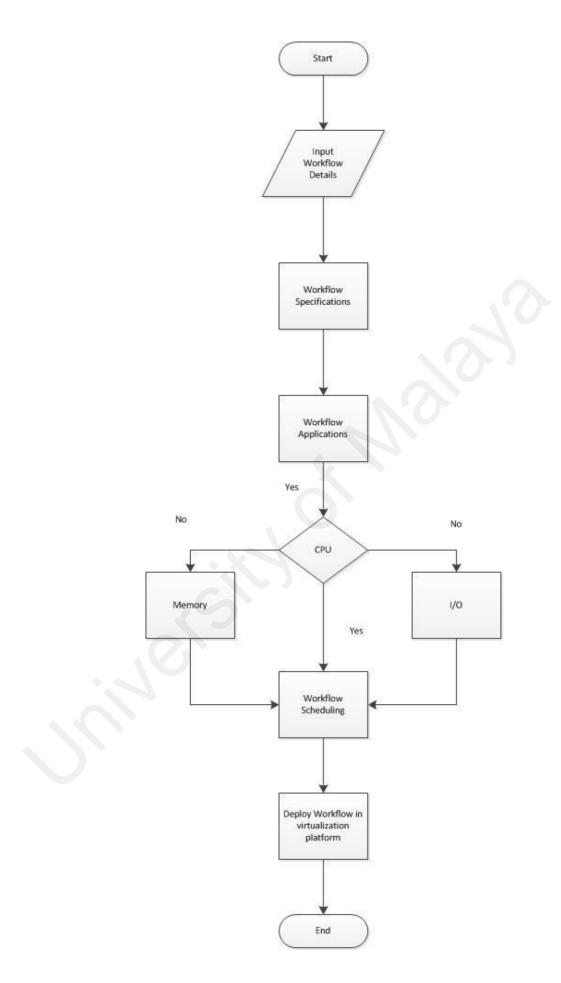


Figure 4.4. System flow

4.2 System Implementation

At this phase, the system is implemented based on the defined automated system architecture. Firstly, a performance database is developed. Then, a HTML form is developed to ease user submission. Finally, the core components, namely orchestration, scheduling and deployment are implemented.

4.2.1 Performance Database

Performance database is developed to gather and store the resource benchmark data. A set of software tools is used to conduct the performance evaluation on the hypervisorbased and container-based platform in terms of compute, memory, and I/O. SysBench and Bonnie++ is selected as the evaluation tool for performance benchmark. Y-Cruncher and STREAM will not be used as there are already a lot of other researchers using it for performance evaluations. Furthermore, SysBench is an effective tool to measure the performance of the applications. SysBench can handle looping test and the limit of each test case can be configured. Bonnie++ is a small utility with the purpose of benchmarking file system IO performance.

4.2.1.1 Compute Performance Testing Using Sysbench

Figure 4.5 shows the compute performance testing. The purpose of this test is to measure how long would it takes for the system to fully execute this CPU intensive test. This test verifies prime numbers with sequentially increasing and verifying that the remainder or modulo calculation is zero. If it is, then the number is not prime and the calculation goes on to the next number. Hence, if none have a remainder of 0, then the

number is prime. The maximum number that it divides by is calculated by taking the integer part of the square root of the number.

root@ubuntu:/home/amin# sysbench - sysbench 0.4.12: multi-threaded s	test=cpucpu-max-prime=3000 run system evaluation benchmark
Running the test with following op Number of threads: 1	otions:
Doing CPU performance benchmark	
Threads started! Done.	
Maximum prime number checked in CF	PU test: 3000
Test execution summary:	
total time:	5.0126s
total number of events:	10000
total time taken by event exec	cution: 4.5537
per-request statistics:	
min:	0.31ms
avg:	0.46ms
max:	9.91ms
approx. 95 percentile:	0.73ms
Threads fairness:	
events (avg/stddev):	10000.0000/0.00
execution time (avg/stddev):	4.5537/0.00

Figure 4.5. Compute performance testing

4.2.1.1.1 Compute Performance Testing Using Bonnie++

Figure 4.6 shows the compute performance testing using Bonniee++. Basically, it will create a test to stress out the CPU and I/O of the system. File size need to declare in the command before executing the test. Normally the size of the file would be double the size of the RAM of its host. After finish executing the test, the result would be saved in the performance database.

amin@ubuntu:~	5 bonni	le++ -d	/tmp -	s 2G	-n 0	-m TES	ST -f	-b -u	amin			
Using uid:1000	D, gid:	1000.										
Writing intel	ligent	.vdon	ie									
Rewritingdo	-	-										
Reading intel	ligent	vdon	ie									
start 'emdo				done.	dor	ne						
Version 1.97							Se	quent	ial Ing	out-	Rand	om-
Concurrency												
		K/sec %										
TEST	2G					946 9			10540			
Latency			12	41ms		327ms			537	IS	975m	s
1.97,1.97,TEST	T,1,151	7392580	,2G,,,	,5387	8,4,4	43046,9	9,,,10	5401,	18,128	1,11		
,,,,,,,,,1241												

Figure 4.6. Compute Performance Testing using Bonniee++

Versi	Version 1.97 Seq			iential Output				Sequential Input				Random	
	Size	Per Cha	r	Block		Rewrit	e	Per Cha	r	Block		Seeks	
		K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	/sec	% CPU
TEST	2G			53878	4	43046	9			105401	18	128.1	11
IESI	Latency			1241m	s	327ms				537ms		975ms	

Table 4.2. CPU Result of Bonnie ++

Table 4.2 show the result of CPU performance test using Bonniee++. There are three columns involved in the result which is the sequential output, sequential input and random seeks. These results are gathered from the test that run on the Linux terminal. The %CPU column reports the percentage of the CPU that was used to perform the IO for each test. The file metadata tests are shown in the second row and files are created, read and finally deleted. The create, read, delete metadata tests are performed using file names that are sorted numerically.

4.2.1.2 Memory Performance Testing

Figure 4.6 shows the memory performance testing. A set of data size is used for the evaluation. The write operation is conducted. A test summary is generated which consists of the number of operations performed, the total time taken and the total number of events. The minimum, average and maximum time are shown as well in the summary.

```
root@ubuntu:/home/amin# sysbench --test=memory --memory-block-size=1M --memory-total-size=100G
--num-threads=1 run
sysbench 0.4.12: multi-threaded system evaluation benchmark
Running the test with following options:
Number of threads: 1
Doing memory operations speed test
Memory block size: 1024K
Memory transfer size: 102400M
Memory operations type: write
Memory scope type: global
Threads started!
Done.
Operations performed: 102400 ( 3339.44 ops/sec)
102400.00 MB transferred (3339.44 MB/sec)
Test execution summary:
     total time:
total number of events:
                                                    30.66385
                                                    102400
     total time taken by event execution: 26.4068
per-request statistics:
           min:
                                                           0.20ms
           avg:
                                                           0.26ms
           max:
                                                          20.72ms
           approx. 95 percentile:
                                                           0.23ms
Threads fairness:
events (avg/stddev):
execution time (avg/stddev):
                                            102400.0000/0.00
                                            26.4068/0.00
```

Figure 4.7. Memory performance testing

4.2.1.3. I/O Performance Testing Using Sysbench

Figure 4.7 shows the I/O performance testing. Files with different size are created, followed by read and write operations. The read and write ratio is defined at 1:50 and a synchronous I/O mode is carried out for the test. Besides, user can specify the block size. A test summary is generated, which consists of the numbers of read and write operations,

speed, and total time. The minimum, average and maximum time are shown as well in

the summary.

```
root@ubuntu:/home/amin# sysbench -·test=fileio -·file-total-size=5G -·file-test-mode=rndrw -
 -rng=on --max-time=300 --max-requests=0 run
sysbench 0.4,12: multi-threaded system evaluation benchmark
 Running the test with following options:
Number of threads: 1
Initializing random number generator from timer.
Extra file open flags: 0

128 files, 40Mb each

5Gb total file size

Block size 16Kb

Number of random requests for random IO: 0

Read/Write ratio for combined random IO test: 1.50

Periodic FSYNC enabled, calling fsync() each 100 requests.

Calling fsync() at the end of test, Enabled.

Using synchronous I/O mode

Doing random r/w test

Threads started!

Time limit exceeded, exiting...

Done.
 Done.
Operations performed: 46582 Read, 31054 Write, 99328 Other = 176964 Total
Read 727.84Mb Written 485.22Mb Total transferred 1.1846Gb (4.0431Mb/sec)
258.76 Requests/sec executed
 Test execution summary:
         total time: 300.0309
total number of events: 77636
total time taken by event execution: 241.6577
per-request statistics:
                                                                                         300.0309s
                   min:
                                                                                                       0.10ms
                                                                                                  3.11ms
256.09ms
                    avg:
                     max:
                    approx. 95 percentile:
                                                                                                     14.08ms
 Threads fairness:
         events (avg/stddev):
execution time (avg/stddev):
                                                                             77636.0000/0.00
                                                                            241.6577/0.00
```

Figure 4.8. I/O Performance testing using Sysbench

Once the performance evaluation is conducted, the benchmark information is stored in the performance database. The database is used as the knowledge information base by the scheduling components.

4.2.1.3.1 I/O Performance Testing Using Bonnie++

Figure 4.9 shows the I/O performance testing using Bonniee++. Basically, the test is the same as the CPU stress test as Bonnie++ can evaluate the CPU and I/O eventually. Hence, for this I/O test, the file size double-up from 2GB to 4GB in order to observe the input and output execution time and size.

	00, gid:1000. lligentlydone						
Rewritingdone							
Reading intelligentlydone							
start 'em	donedonedon	edone	.done				
	7Seq						
Concurrency	1 -Per Chr-	Block	-Rewrite-	-Per Chr	-Block	Seeks	
Machine	Size K/sec %CP	K/sec %CP	K/sec %CP	K/sec %CP K	/sec %CP	/sec %CP	
TEST	4G	57085 4	41878 10	1	00943 18	151.1 14	
Latency		800ms	404ms		535ms	870ms	
and a second							

Figure 4.9 I/O Performance Testing Using Bonnie++

V	ersi	on 1.97	1.97 Seq			uential Output			Sequential Input					
		Size	Per	Char	Block		Rewrit	e	Per (Char	Block		See	eks
			K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	K/sec	% CPU	/sec	% CPU
т	EST	4G			57085	4	41878	10			100943	18	151.1	14
11	201	Latency			800ms		404ms				535ms		870ms	

 Table 4.3. I/O Result of Bonnie ++

Table 4.3 show the result of I/O performance test using Bonniee++. There are three columns involved in the result which is the sequential output, sequential input and random seeks. These results are gathered from the test that run on the Linux terminal. The

sequential output will show the kilobyte of block per second were involved in output portion. The sequential input shows the kilobyte of block per second involve in input portion.

4.2.2 Web Form

A Web form is developed to ease workflow submission. The proposed system is developed by using the Java programming language. Figure 4.8 shows the Web form that allows user to input the workflow names, as well as the specifications. User is required to input in XML file as well as all the executable workflow jobs in the workflow.

lobeo Fill in bolow detaile and rink	1 10 11 11	
lease Fill III below details and upi	oad workflow applications	
Workflow Name		
Workflow Specification, Plea	se upload the XML File	
Choose File No file chosen		
Upload Application		
Choose File No file chosen		

Figure 4.10. Web form

Figure 4.9 shows the workflow specifications that can be uploaded into the automated system architecture. Each workflow has multiple subtasks. The subtask is defined by the subtask ID. Each subtask contains the task types that are divided into compute, memory

and I/O. The execution command needs to be clarified on every task in order for the job to be successfully executed. After the execution is completed, the system output will be automatically saved into the respective directory path folder.

```
<job username="amin" directory="/star/u/amin/test" title="My test program" description="Testing my program">
<task number="4">
<subtask id="1" tasktype="C" value="10000">
<command>myProgramCompute -c -d</command>
<stdin URL="nfs:/star/u/amin/test/input1" />
<stdout URL="nfs:/star/u/amin/test/output1" />
</subtask>
<subtask id="2" tasktype="M" value="10240">
<command>myProgramMemory -c -d myFile</command>
<stdin URL="nfs:/star/u/amin/test/input2" />
<stdout URL="nfs:/star/u/amin/test/output2" />
</subtask>
<subtask id="3" tasktype="I" value="200">
<command>myProgramIO -c -d myFile</command>
<stdin URL="nfs:/star/u/amin/test/input3" />
<stdout URL="nfs:/star/u/amin/test/output3" />
</subtask>
<subtask id="4" tasktype="C" value="10000">
<command>myProgramCompute -c -d</command>
<stdin URL="nfs:/star/u/amin/test/inpu4t" />
<stdout URL="nfs:/star/u/amin/test/output4" />
</subtask>
</task>
</job>
```

Figure 4.11. Workflow specifications

4.2.3 Orchestration Component

Total of Workflow in each Category------ X100%

Total Number of Workflow

This orchestration component will be involved in the evaluation of the task itself. The evaluation of the task will depend on its intensiveness. The intensiveness of every workflow is determined by the total number of category workflow divided by the total number of workflow times 100 percent. Table 4.3 show the example on how to declare the threshold for workflow intensiveness based on their percentage.

Sequence of task in a workflow	Category	Number of workflows on each category	Total number of workflows for each category	Total Percentage for each category
1	CPU Intensive (A)	1/5		
2	CPU Intensive (B)	1/5	3/5	60%
3	CPU Intensive (C)	1/5		
4	I/O Intensive	1/5	1/5	20%
5	Memory Intensive	1/5	1/5	20%
	Total	5/5	5/5	100%

 Table 4.4 Example of Threshold for Choosing Workflow Intensiveness

Hence, the highest percent of the category will be declared as the most intensive category for the workflow. If the task requires a lot of CPU resource allocation, then it will be orchestrated together with another task which is in the same category. If the task requires a lot of memory resources, then it will be separated into the same category. As if the task is I/O intensive, it will be merged with the other task that requires massive I/O resource allocations.

4.2.4 Scheduling Component

After the entire task was orchestrated into their category, the scheduling component will schedule the task into different virtualization platforms, namely hypervisor-based virtualization or container-based virtualization.

4.2.5 Deployment Component

Once all tasks are orchestrated and scheduled, the last component would be the deployment. All tasks and jobs will be deployed to the respective virtualization platform to be executed.

4.3 System Testing

As the purpose of testing the system, a workflow with four tasks is submitted to the system. The first task was a compute intensive job. This job contains many tasks that require compute resources in order for it to be executed. The second task is the memory intensive job, the third task is I/O intensive job and the final task is also compute intensive job.

No	Task Name	Exec Time for Hypervisor	Exec Time for Container	Exec Time for Proposed Architecture
1	Compute Intensive Task (A)	4.55 sec	3.37 sec	3.37 sec
2	Memory Intensive Task	3.06 sec	1.61 sec	1.61 sec
3	I/O Intensive Task	2.90 sec	6.67 sec	2.90 sec
4	Compute Intensive Task (B)	13.65 sec	10.11 sec	10.11 sec
	Total Execution Time	24.16 sec	21.76 sec	17.99 sec

 Table 4.5. Time Execution Comparison Between Workflow

Table 4.2 shows the experimental test results. The table shows the total execution time of running the workflow in hypervisor, container and the proposed architecture.

4.4 Chapter Summary

In this chapter, the system design of the proposed system architecture is described in detail. Then, the system implementation is shown and explained. In the last phase, the system testing procedures are discussed. In the next chapter the experiment testing is discussed followed by result analyst and discussion.

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CHAPTER 5: RESULTS AND DISCUSSION

This chapter details the performance evaluation process of the proposed technique and discussions on the obtained results. The chapter also presents the experimental setup by using the hypervisor-based and container-based virtualization. The tests are carried out to capture the result, which is the total execution time for each virtualization platform. After the test is completed, the obtained results will be analysed for each technique. A chapter summary is presented at the end of chapter.

5.1 Experimental Setup

The setup of the performance evaluation test is to evaluate the time taken for each hypervisor-based and container-based virtualization platform to execute respective jobs and tasks. Four different types of workflow specification are used, namely compute intensive, memory intensive, I/O intensive and uniform intensive. Compute intensive workflow consists of a higher rate of compute intensive tasks as per shown on Table 4.4, which there are more than 60% of compute intensive task occur in one workflow. Same goes with the memory intensive workflow which it consist a higher rate of memory intensive tasks while I/O workflow consists of a higher rate of a higher rate of I/O intensive tasks. The uniform workflow is the mixture of compute, memory and I/O intensive tasks. The uniform workflow consists of a nearly equal distribution of the compute, memory and I/O intensive tasks.

Table 5.1 shows the example of some of the workflow that was executed from the test case. As shown in the table, every workflow has their own task and this task was sorted by sequence based on their priority. Each workflow consisted of multiple tasks. Each task was categorised as compute, memory or I/O intensiveness task. After the task and its

value are defined, all task will be deployed, and the execution time will be captured accordingly. The total execution time will be analysed and documented.

Workflow Number	Task Name	Value	Total time (sec)
1	Compute (A)	10000	4.55
	Memory (A)	10240	3.06
	I/O (A)	200	2.90
	Compute (B)	10000	4.55
	I/O (B)	200	2.90
	Memory (B)	10240	3.06
2	Compute (A)	20000	10.05
	Compute (B)	10000	4.55
	Compute (C)	5000	2.28
	Memory (A)	20480	6.12
	Memory (B)	10240	3.06
3	Compute (A)	20000	10.05
	Compute (B)	10000	4.55
	Memory (A)	20480	6.12
	Memory (B)	10240	3.06
	Memory (C)	40960	12.24

 Table 5.1. Workflow specifications

Table 5.2 shows the hardware specifications that are used for the experimental testing. Intel Pentium G2010 @ 2.80 GHz is used as the processor, with 10GB physical memory. The PC was running on Linux Ubuntu Version 14.04 64-bit.

Table 5.2. Host specifications

Host Specifications					
Processor	Intel Pentium G2010 @ 2.80 GHz				
RAM	10 GB				
System Type	64-bit Operating System				
Operating System	Linux Ubuntu 14.04				

5.2 Experimental Results

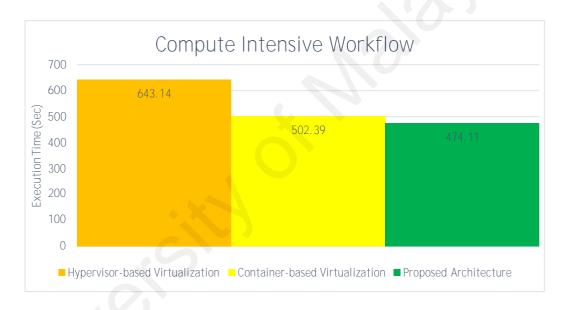


Figure 5.1 Compute intensive workflow execution time result

Figure 5.1 shows the results and output of the experiment task tested on the compute intensive workflows. The X-axis represents the type of virtualization and the proposed architecture while the Y-axis represents the total execution time. As shown in Figure 5.1, the execution time for the proposed architecture is 474.11 second which is the fastest. When the workflow is executed on container-based virtualization, it takes 502.39 second, which is 5.96% more than the proposed architecture. The workflow takes the longest time

to execute on hypervisor-based virtualization which is 35% longer than the proposed architecture. The compute intensive workflow consists of higher percentage of computation tasks. Since the proposed architecture can select the best virtualization to execute the task, the total execution time is minimised as compared to hypervisor-based and container-based virtualization.

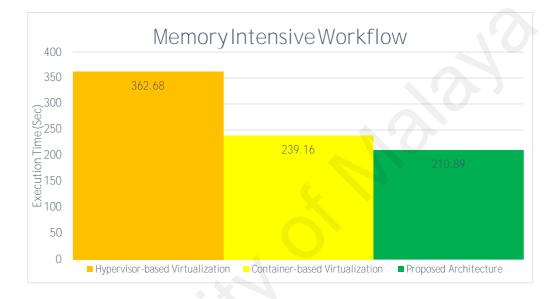


Figure 5.2. Memory intensive workflow execution time result

Figure 5.2 shows the results and outputs of the experiment testing on the memory intensive workflows. The X-axis represents the type of virtualization and the proposed architecture while the Y-axis represents the total execution time. It takes 210.89 second to execute the workflow by using the proposed architecture. The workflow takes 362.68 second and 239.16 second to complete when running on hypervisor-based and container-based virtualization respectively. The result shows that the proposed architecture minimises the total execution time as compared to hypervisor-based and container-based virtualization.

Figure 5.3 shows the results of I/O intensive workflow test for the hypervisor-based, container-based and proposed architecture. The workflow takes 255.61 sec to complete when running using the proposed architecture. As for the hypervisor-based virtualization, it takes 287.72 sec to complete, which was 12.56% longer as compared to the proposed architecture. The workflow takes the most time which 516.87 sec is to complete when running on container-based virtualization.

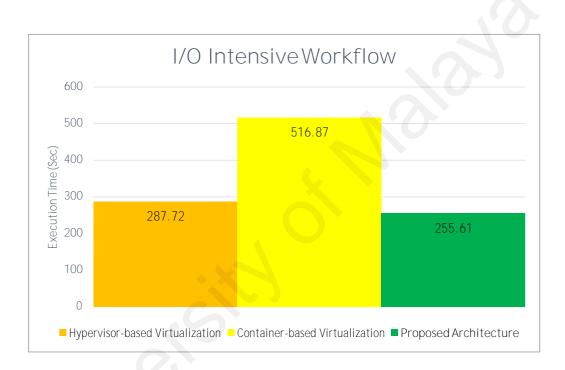


Figure 5.3. I/O Intensive workflow execution time result

The I/O intensive workflow consists of higher percentage of I/O tasks. These tasks are basically to read and write into the disk or data store. I/O-intensive environments can have large, medium or small amounts of storage, yet have active workloads with either many I/O (transactions) of various sizes. The result shows that the proposed architecture minimises the total execution time as compared to hypervisor-based and container-based virtualization.

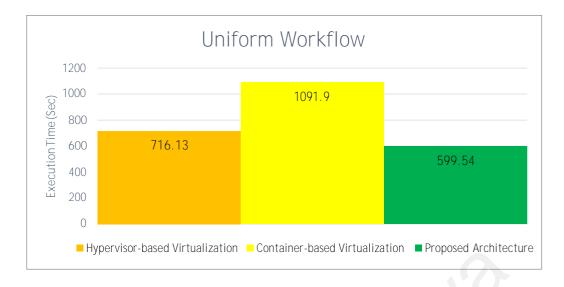


Figure 5.4 Uniform Intensive Workflow Execution Time Result

Figure 5.4 shows the result of a uniform workflow testing. A uniform workflow consists of a combination of compute, memory and I/O intensive tasks. As shown in Figure 5.4, the total execution time was 599.54, 716.13 and 1091.9 second when running on proposed architecture, hypervisor-based and container-based virtualization respectively. The total execution time is 19.44% and 82.12% longer when running on hypervisor-based and container-based virtualization respectively, as compared to proposed architecture. The result shows that the proposed architecture minimises the total execution time as compared to hypervisor-based and container-based virtualization.

5.2 Chapter Summary

In this chapter the workflow testing of the hypervisor-based, container-based and proposed architecture was presented. Based on the experimental results, the proposed architecture minimises the execution time for all workflows as compared to hypervisorbased and container-based virtualization.

CHAPTER 6: CONCLUSION AND FUTURE DIRECTIONS

This chapter summarises the conducted research. It discusses the research contribution in brief. Then, the suggestions for possible improvements are provided for future research plans.

6.1 Thesis Summary

This research starts with exploring the underlying concepts, characteristics, system architecture of the most famous virtualization technology, namely hypervisor-based and container-based. Then, the performance evaluation strategies and tools are discussed and analysed. Subsequently, the scientific workflow is explained briefly, and problem of the current virtualization platform is identified.

During the investigation and analysis process, the problems of the workflows deployment in the virtualization platform are identified. It is very challenging to deploy the workflow tasks to the suitable virtualization platforms since the workflow consists of compute, memory and I/O intensive tasks.

To address the above issue, an automated system architecture is proposed. The proposed architecture consists of the core components namely orchestration, scheduling and deployment that are able to minimise the workflow total execution time.

To validate the accuracy and effectiveness of the proposed architecture, the benchmark is conducted to measure the performance of the virtualization platform. The performance evaluation is conducted by using the performance evaluation tools called SysBench. Once the result is obtained, the information is used as the knowledge-based for the proposed automated architecture. Four different types of workflows are used for experimental testing. The experimental result indicates that the proposed automated architecture outperformed the standard deployment of the workflow in the hypervisor-based and container-based virtualization platform by minimising the total execution time approximately 20%-30%. It is proven that the proposed automated architecture reduces the total execution time as compared to hypervisor-based and container-based virtualization.

6.2 Thesis Contribution

In this research, an automated system architecture is proposed and implemented. The workflow orchestration, scheduling and deployment components are developed by using Java programming language. Then, the proposed architecture is evaluated and tested by using different workflow tasks, namely compute, memory and I/O intensive tasks. The analysis is done based on the result that gained from the SysBench performance testing tools. The experimental results indicated that the proposed architecture can minimise the total execution time.

To summarise, this research has contributed in few areas.

First, it provides a deep understanding of workflow orchestration, scheduling and deployment in the automation architecture. The automation system will orchestrate, group and schedule the workflow before being deployed and executed into the respective virtualization platform.

Secondly, in this study, the experimental results show that the proposed automation architecture can reduce the total execution time of the workflow. The effectiveness of the architecture is proven based on the measurements.

Lastly, based on the experimental results, the proposed architecture can be used to manage and deploy workflow in the virtualization environment, either it was on hypervisor-based or the container-based.

6.3 Future Work Suggestions

Many research was done in hypervisor-based and container-based technologies. However, there are many issues and challenges which have directed the research to be further studied and explored. One of the major issues is the application performance on the virtualization environment as many have claimed that their applications performance was decreased after it was migrated into the virtualization platform.

A suggestion to further improve the current research is to consider the network bandwidth of both virtualization platform since a network plays a big role in managing applications. In addition, it is suggested to include the size of the workflow that will be deployed to see if there is any difference if a mass size of workflow is deployed into both virtualization platforms.

In summary, virtualization is one of the most advanced and emerging technologies for Cloud computing and data centre. The advances in this technology will lead to a broad and deep research in which the outcomes will benefit both administrators and users.

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